Process, Systems and Tests: Three Layers in Concurrent Computation (Short Paper)

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Framework

1. A syntax and a semantics $P \to P'$
2. Some observations $P\downarrow$
3. An order[a]: $Q \leq Q'$ if for every context $C[\,]$ if $C[Q] \to\downarrow$ then $C[Q'] \to\downarrow$

[a] With nondeterminism, play the (bi)-simulation game.

Motto

"*Two terms are equivalent if they produce the same observations in every context.*"

But. . .

who does the context represent?

Contextual Reasoning: Who Is the Context?

A context can represent either. . .

- a (hostile) environment,
- some other part of the program.

Not the same point of view!

Let's take the perspective of a programming language, and see the difference between a snippet of code and a function.

A snippet

```
while(i < 10){
    System.out.print(i);
    i ++;
}
```

A snippet

```
public class HelloWorld{
    public static void main(String []args){
        int i = 0;
        while(i < 10){
            System.out.print(i);
            i ++;
        }
    }
}
```

What can the context do?

Add more instructions, declare the variables, declare the function name, . . .

A snippet

```
public class HelloWorld{
    public static void main(String []args){
        int i = 0;
        for (; i<9; i++){
            System.out.print(i+1);

        }
    }
}
```

Who is the context?

The context is the programmer. Two snippets are the same if the programmer can subtitute one for the other during programming.

A Function

```
public class Toto{
  public int f(int x){
    int i = 0; int a = 0;
    while(i < x){
      a = a+i;
      i ++;
    }
    return a;
  }}
```

What can the context do?

Call the function!

A Function

```
public class Toto{
  public int f(int x){
    int i = 0; int a = 0;
    while(i < x){
      a = a+i;
      i ++;
    }
    return a;
  }}
```

```
public class HelloWorld{
  public static void main(){
    Toto t = new Toto();
    System.out.print(
        t.f(12)
    );
  }
}
```

Who is the context?

The context is the user (potentially hostile). Two functions are the same if the user cannot distinguish them by calling them.

Usual Syntax

$$P = 0 \mid a(x).P \mid \bar{a}(M).P \mid (\nu a)P \mid P \parallel P \mid P + P \mid \cdots$$

Which contexts can we use?

Programmer "All!"

"I'm *building* my process."

User "Only parallel composition."

"I'm *communicating* with my process."

Yet, contexts are *on the surface* often defined as a monolithic concept.

On Open and Closed Terms

Open terms Still being built.

Closed terms Ready to be released.

The original π calculus presentation blurs the distinction ("Everything is a name").
Later presentations reintroduce it.

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Our Proposal: Processes, Systems and Tests

1. Define processes.

2. Define deployment criteria to transform processes into systems.

3. Define tests on systems.

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   1. Choose the *construction operators*.
   2. The programmer use them to write *processes*.
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1. Define processes.

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   1. Binding of variables
   2. The construction operators at top-level

   3. *Deployment operators*

   A system is *complete*.
3. Define tests on systems.

Our Proposal: Processes, Systems and Tests

1. Define processes.

2. Define deployment criteria to transform processes into systems.
   1. Binding of variables (= restriction)
   2. The construction operators at top-level (= no unguarded sum at top level)
   3. *Deployment operators* (it could be more general, and restricts, expands or intersects w/ construction operators)

   A system is *complete*.
3. Define tests on systems.

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1. Define processes.

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3. Define tests on systems.
   1. Observables, i.e. functions from systems to a subset of a set of atomic proposition (=“emits barb *a*”, “terminates”, “contains recursion operator”, etc.),
   2. A notion of context, that should come with its own set of *testing operators* and reduction rules.

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   For instance: traditionally in π one may consider only *static contexts*.

(Pre-) Congruences

are defined thanks to test as

"observationally,"

"contextually-closed,"

"reduction-closed[a]",

relations.

[a] In presence of nondeterminism

Two (of the) Perks

1. Every step allows the introduction of operators,
2. Multiple notions of systems or tests can (and should!) co-exist in the same process algebra, one being targeted to e.g. programmers, and another to e.g. users.

No need to decide which set of observations is more basic.

(Usual) Context Lemma

Given

fixed sets of processes, systems and observations,

two sets of testing contexts $S \subseteq S''$

the observational congruences defined by the sets of context S and S' coincide.

Examples

In the π-calculus the context $a\langle M \rangle.[\,]$ does not change the congruence,

In the λ-calculus applicative contexts suffice.

Can we define a metatheorem?