

Perfect Sampling of Load Sharing Policies in Large Scale Distributed Systems

Gaël Gorgo and Jean-Marc Vincent

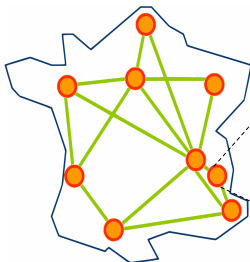
MESCAL-INRIA Project
Laboratoire d'Informatique de Grenoble
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Checkbound meeting
21 Octobre 2010



Large scale computing

Grid 5000

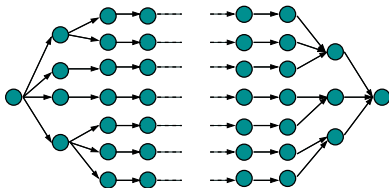


Grenoble cluster



300 cores

Workload model



Load sharing middleware

- Distributed control algorithm
- migration of tasks between nodes

Practical needs

Load sharing policy

A controller (local) checks the utilization of the node and decides to share some work with other nodes.

- When? → Control triggering
- Who? → Paradigm Push, Pull
- Decide? → Local state condition
- How many? → Amount of work to be transferred
- Where? → Selection among targets (probing scheme)

Users requirements

- Maximize the utilization of resources (number of active nodes)
- Minimize the network utilization (number of transfers, costly transfers)

⇒ Need of a tool to evaluate the load sharing policies



Performance evaluation of Load sharing systems

Methodology

- 1 Quantification of the system : **steady-state evaluation**
- 2 Comparison of systems, paradigms, policies
- 3 Tuning of system parameters

Numerical approaches

- Markovian modelling and direct numerical solving
- Matrix geometric solution [ELZ86, MTS90]
- Mean field [Mit98, BGY98]
- Simulation [KH02, DKL98]

Key challenge : very large state space (C^K)



Steady-state simulation of Markov models

Generate typical state, i.e. distributed according to the steady-state

forward simulation

Run from an initial state and stop after a sufficiently long period
⇒ Choice of a stopping rule

Perfect simulation [PW96]

Coupling from the past scheme

- Exact stopping criteria
- Unbiased sampling
- **Monotonicity implies simulation efficiency**

Are the load sharing systems monotone so that we can simulate them efficiently?



Outline

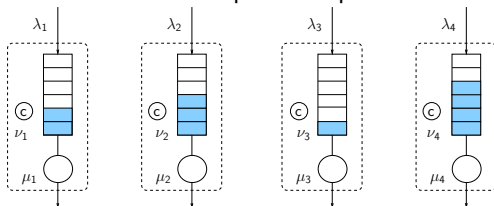
- 1 Large scale systems evaluation
- 2 Modelling of Load sharing systems
- 3 Monotonicity of control policies
- 4 Applications
- 5 Conclusion

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Load sharing model

Parallel independent queues



State space : number of tasks in each queue; $\mathcal{X}_1 \times \dots \times \mathcal{X}_K$

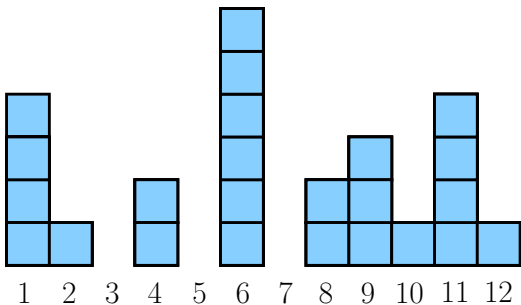
Dynamics : events driven by Poisson process (Poisson system [Bre99]) :

- Generation of a new task in a queue, with rate λ
- Task completion, with rate μ
- Control, with rate ν

Uniformization \Rightarrow Stochastic Recurrence Equation $X_{n+1} = \Phi(X_n, E_{n+1})$

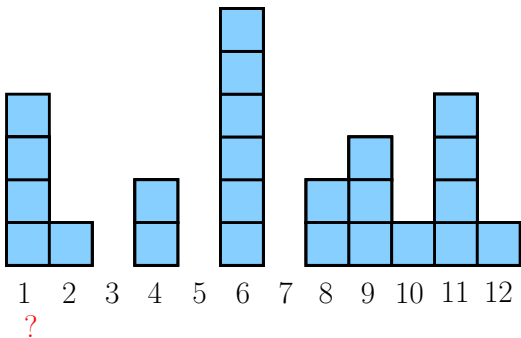
Control event example 1 : *PSQ*

Push to the least loaded node among potential targets (Push to the Shortest Queue)



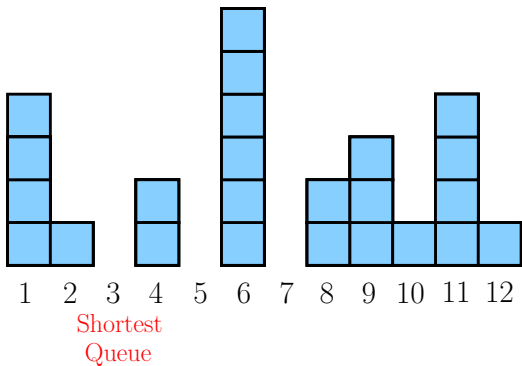
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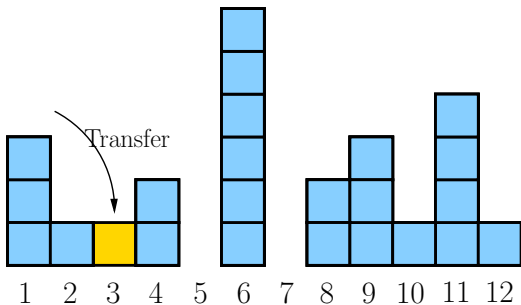
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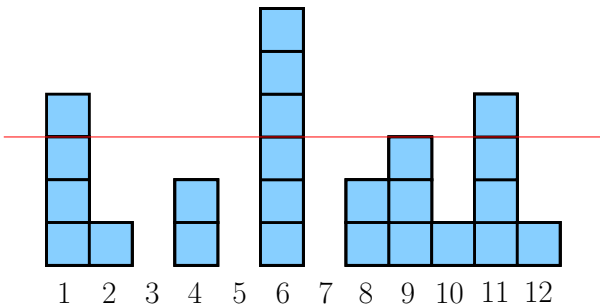
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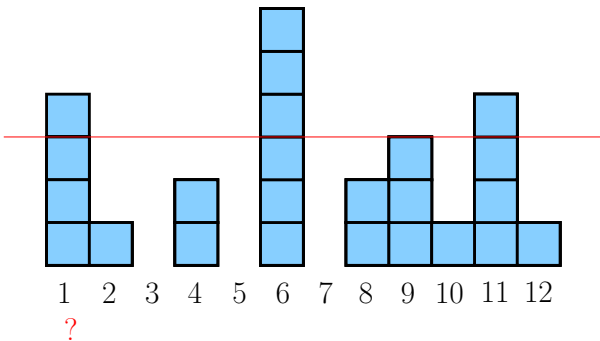
Control event example 2 : *conditionned PSQ*

PSQ Adding an overload threshold on the origin



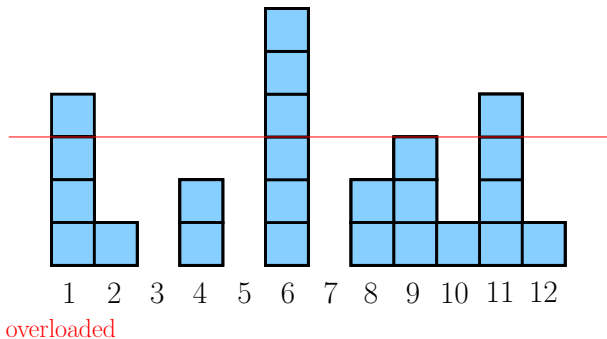
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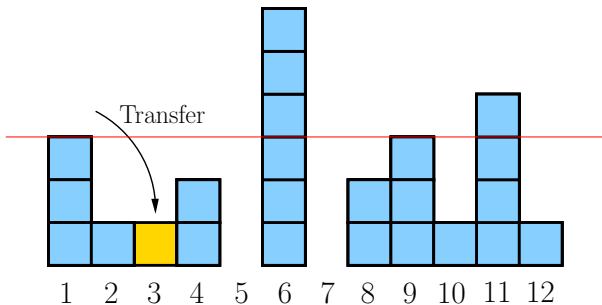
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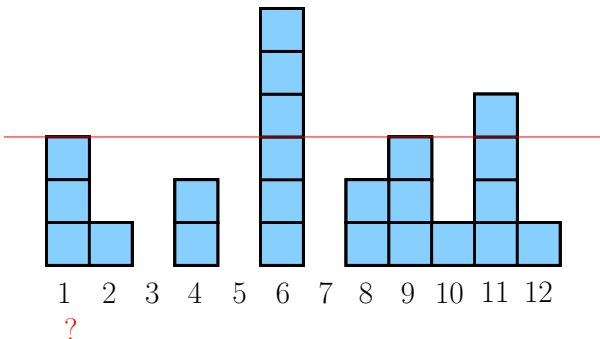
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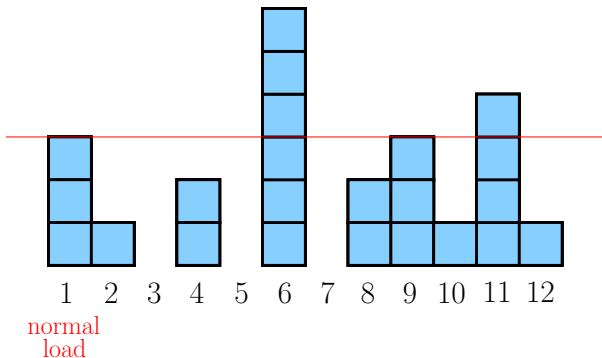
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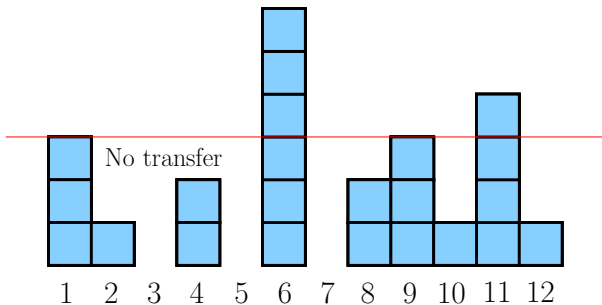
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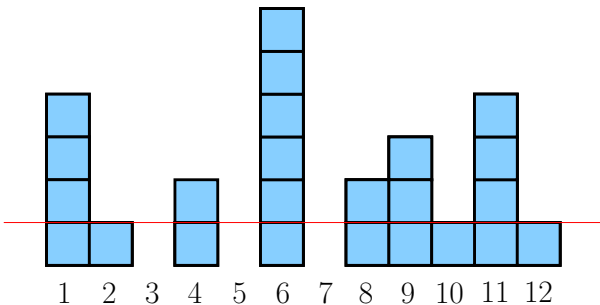
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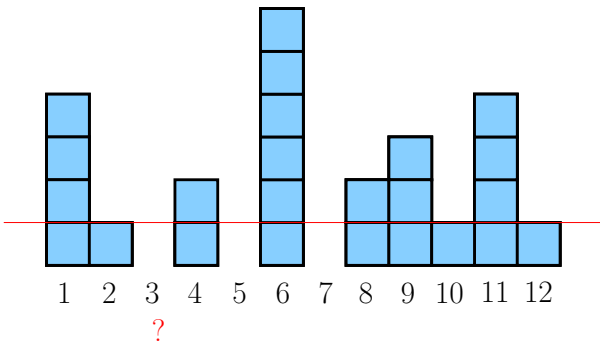
Control event example 3 : *Pull from a probed node*

Pull with probing according to a priority list



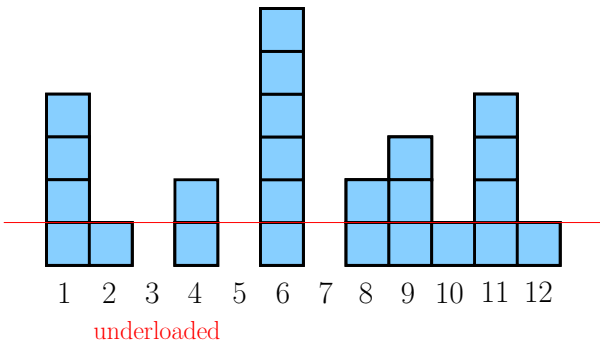
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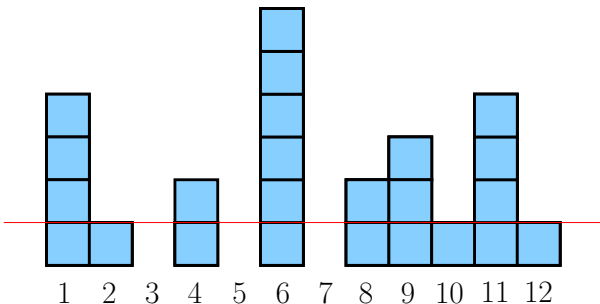
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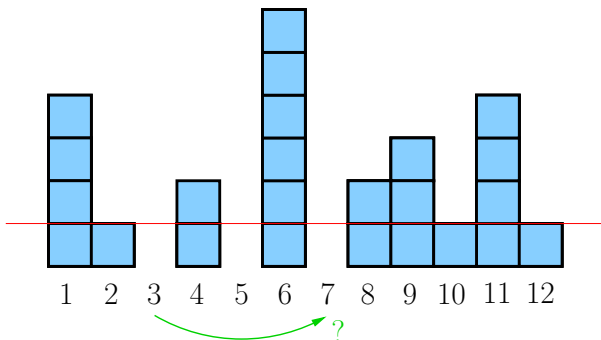
asking potential victims



Prob-limit

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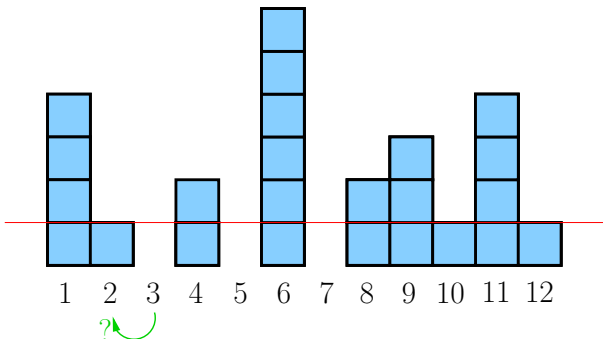
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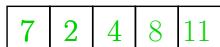
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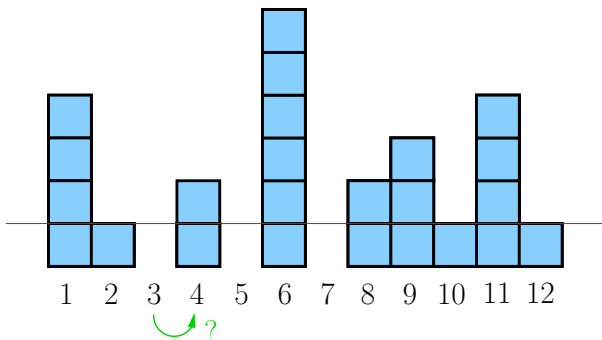


asking potential victims



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asking potential victims



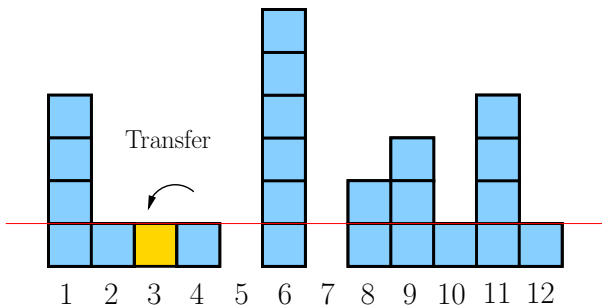
?



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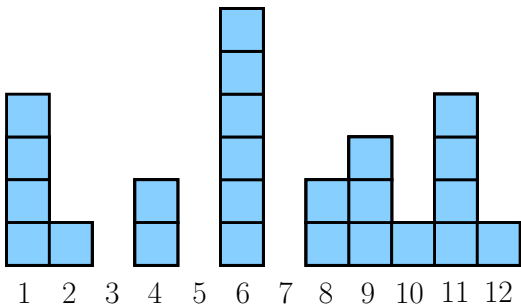
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←—————→
Prob-limit

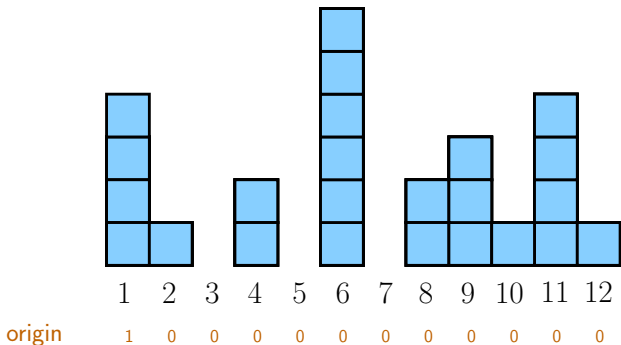
Index model for PSQ

Push to the least loaded node among potential targets (Push to the Shortest Queue)



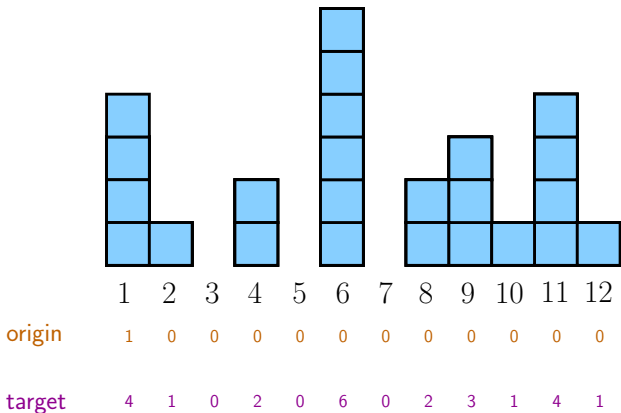
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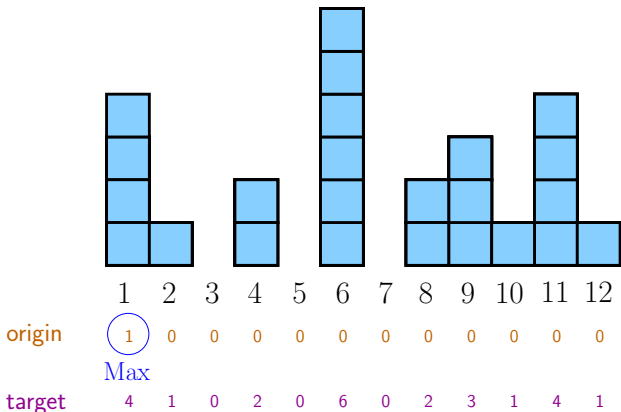
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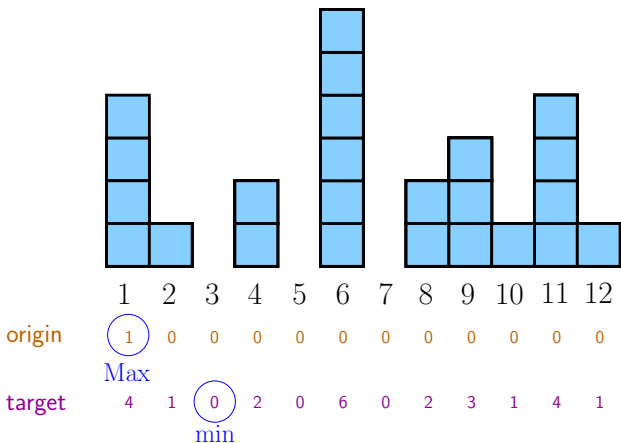
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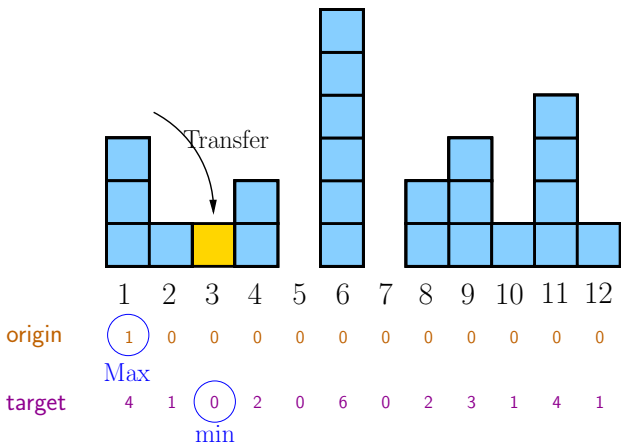
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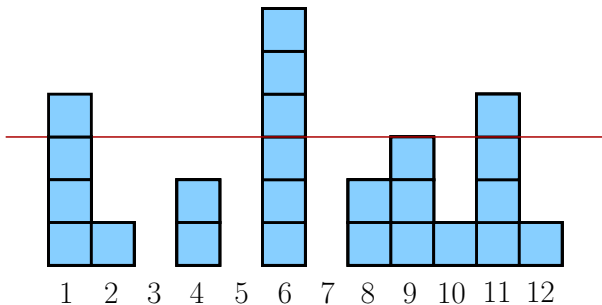
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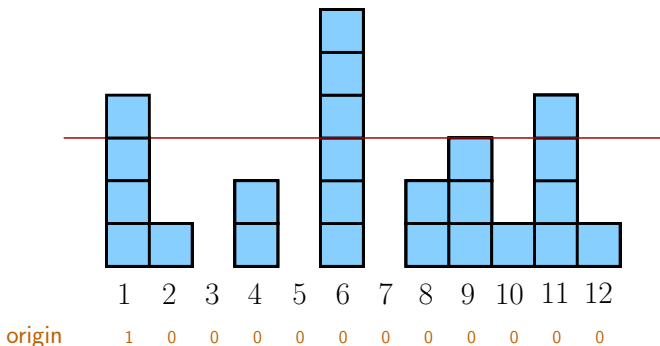
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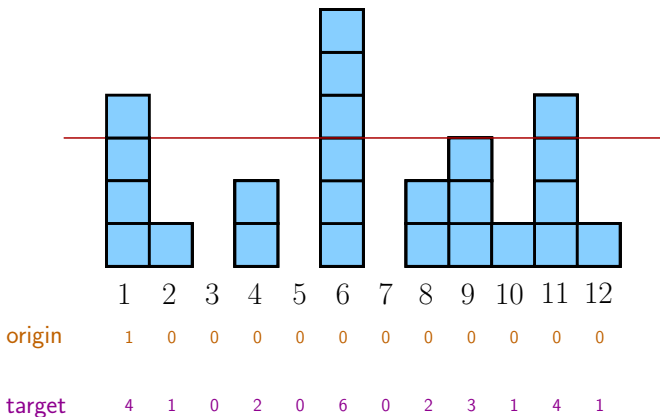
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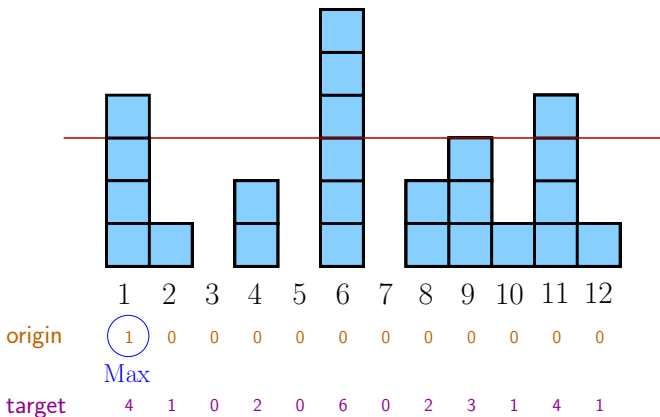
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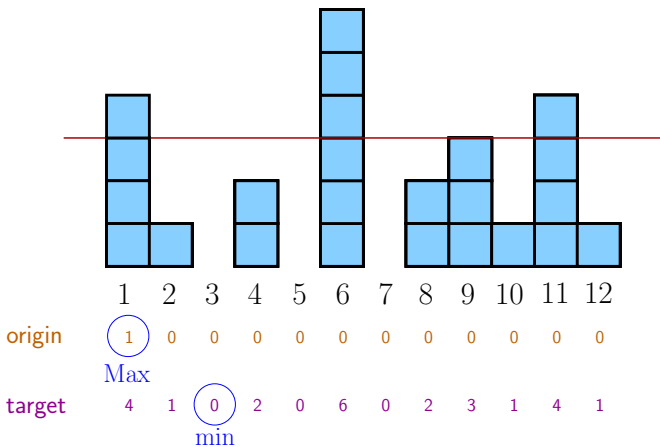
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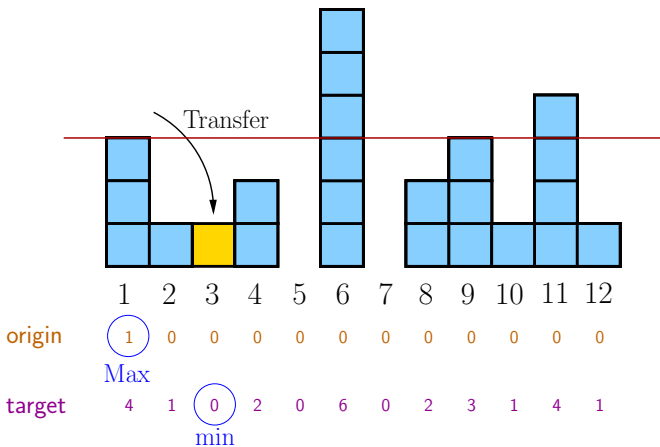
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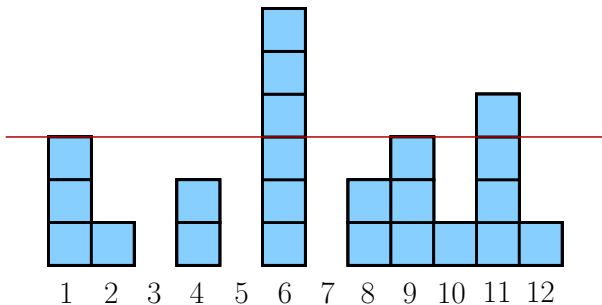
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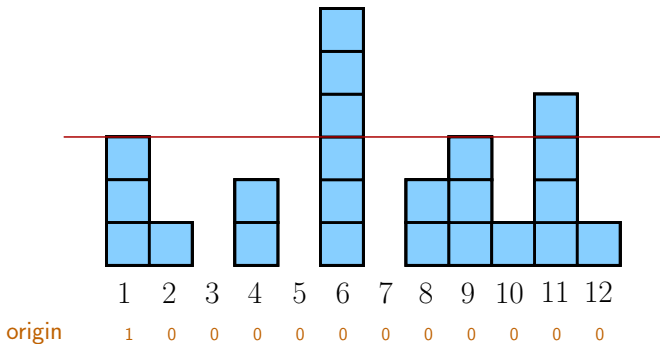
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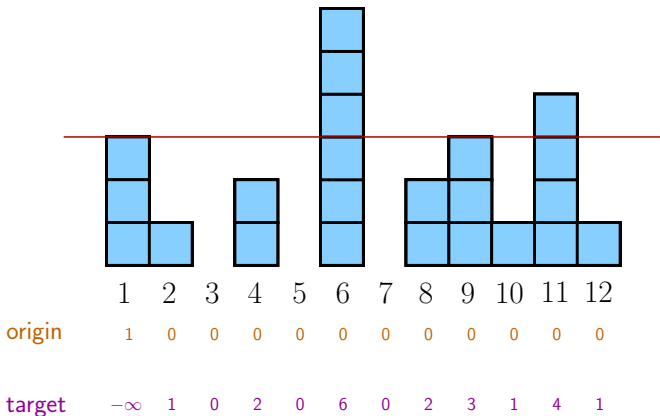
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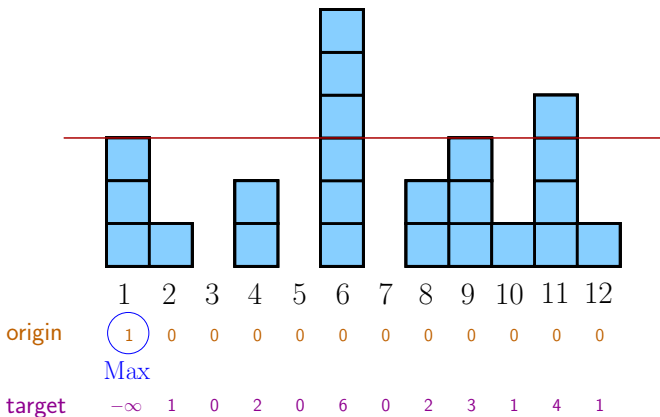
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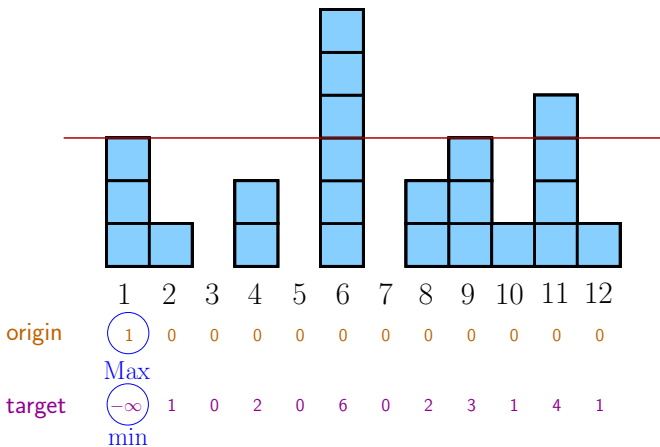
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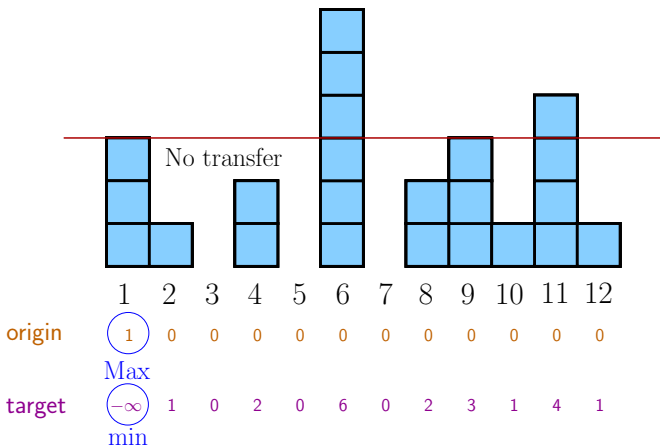
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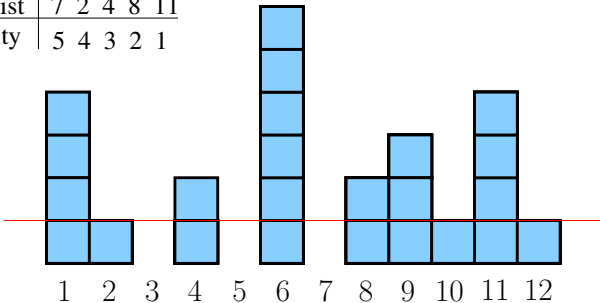
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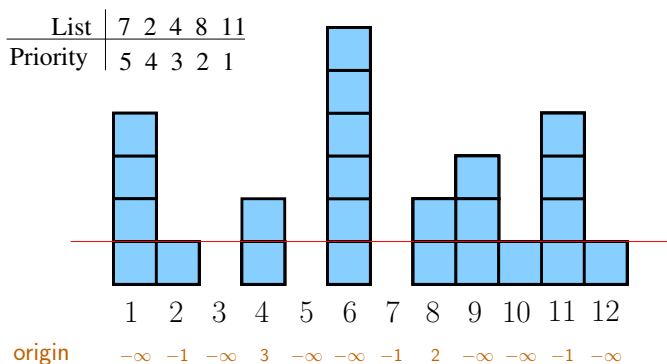
Pull with probing according to a priority list

List	7	2	4	8	11
Priority	5	4	3	2	1



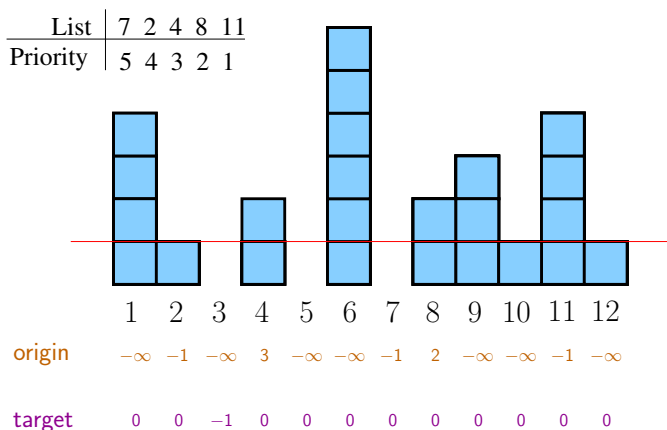
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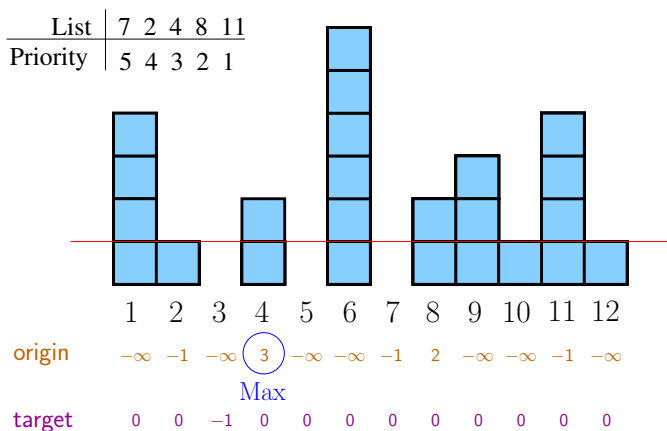
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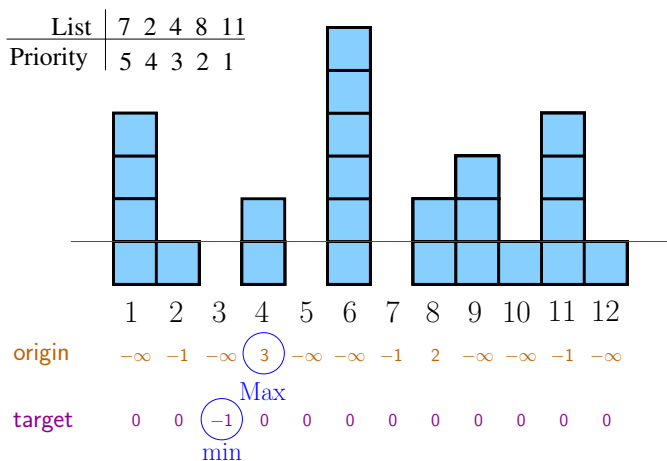
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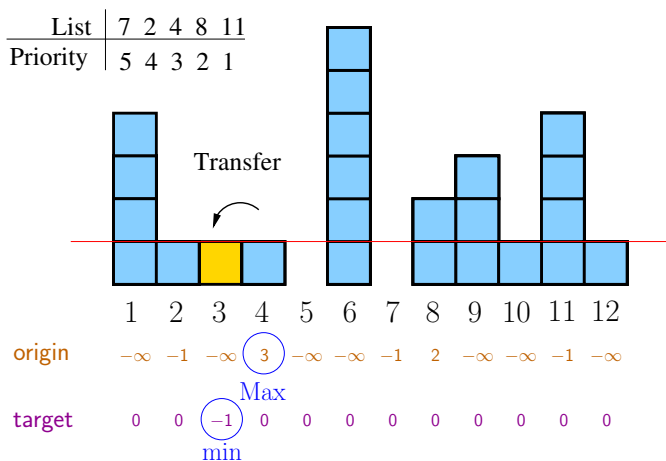
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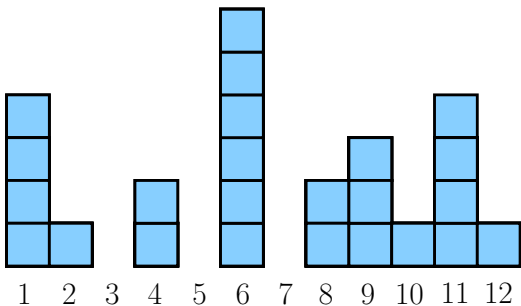
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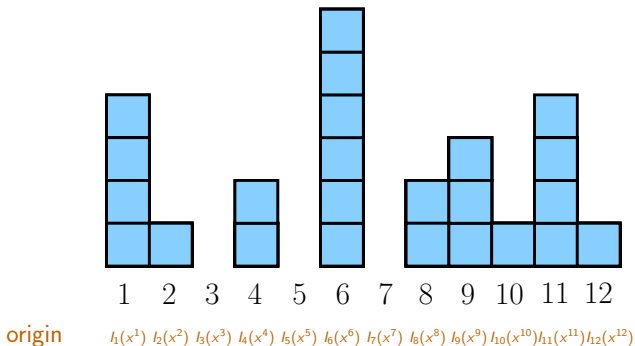
General index model

transfer from an origin (max) to a target (min)



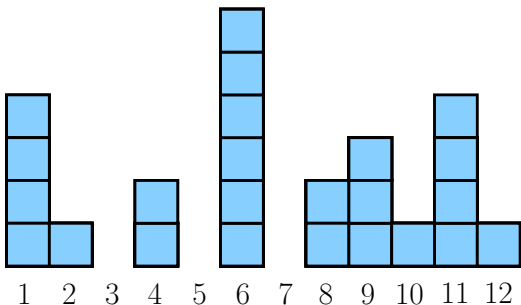
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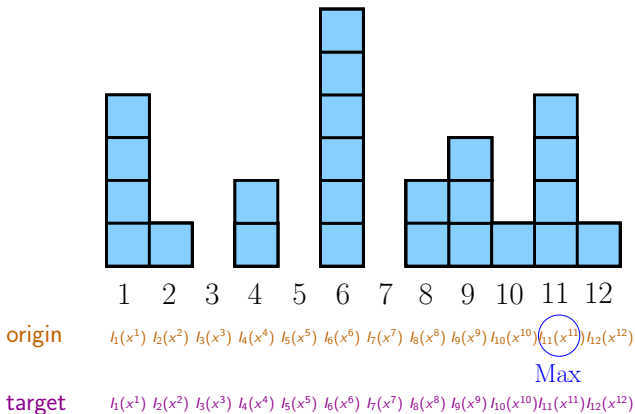
$h_1(x^1) h_2(x^2) h_3(x^3) h_4(x^4) h_5(x^5) h_6(x^6) h_7(x^7) h_8(x^8) h_9(x^9) h_{10}(x^{10}) h_{11}(x^{11}) h_{12}(x^{12})$

target

$h_1(x^1) h_2(x^2) h_3(x^3) h_4(x^4) h_5(x^5) h_6(x^6) h_7(x^7) h_8(x^8) h_9(x^9) h_{10}(x^{10}) h_{11}(x^{11}) h_{12}(x^{12})$

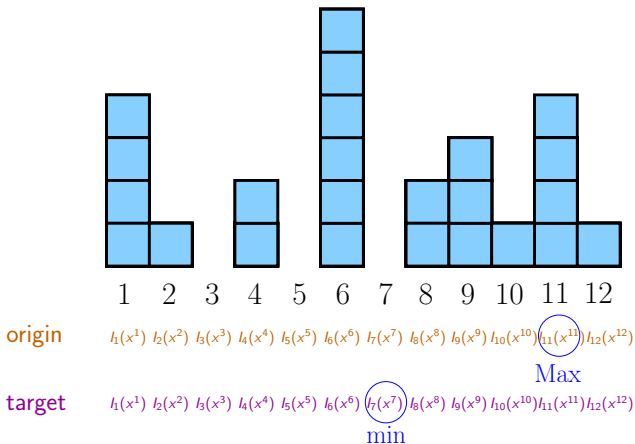
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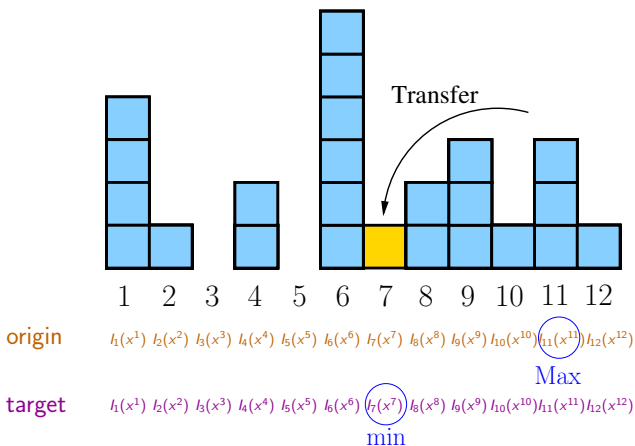
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General index model

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Formalisation

A control event c is defined by :

$$\Phi(x, c) = x - \delta_i + \delta_j$$

i is the **origin**

j is the **target**

Index function

A function $I_k(x^k)$ gives an index, i.e. a cost value to Q_k .

$$i = \mathbf{argmax}_{1 \leq k \leq K} (I_k^{c,o}(x^k))$$

$$j = \mathbf{argmin}_{1 \leq k \leq K} (I_k^{c,t}(x^k))$$

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Monotonicity of index load sharing policies

Monotonicity

- \preceq is the natural partial order on the multi-dimensional state space $\mathcal{X} = \mathcal{X}_1 \times \dots \times \mathcal{X}_K$.

$$x \preceq y \Leftrightarrow x^i \leq y^i \quad \forall i$$

- An event e is monotone if it preserves the partial ordering \preceq on \mathcal{X}

$$\forall (x, y) \in \mathcal{X} \quad x \preceq y \Rightarrow \Phi(x, e) \preceq \Phi(y, e)$$

Theorem

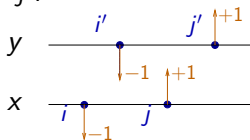
If all index functions $I_k^{c,o}(x^k)$ and $I_k^{c,t}(x^k)$ are monotone and increasing in function of x^k , then the event c is monotone

Proof

Let $x, y \in \mathcal{X}$ two states with $x \preceq y$,

c a control event, $\Phi(x, c) = x - \delta_i + \delta_j$, $\Phi(y, c) = y - \delta_{i'} + \delta_{j'}$.

Suppose that $i \neq i' \neq j \neq j'$.



Then,

$$I_j^{c,t}(x^j) < I_{j'}^{c,t}(x^{j'})$$

$$I_{j'}^{c,t}(x^{j'}) \leq I_{j'}^{c,t}(y^{j'})$$

$$I_{j'}^{c,t}(y^{j'}) < I_j^{c,t}(y^j)$$

$$I_j^{c,t}(x^j) < I_j^{c,t}(y^j)$$

$$x^j < y^j$$

j is the argmin for x
 $I_{j'}^{c,t}$ increasing and $x^{j'} \leq y^{j'}$

j' is the argmin for y

by transitivity

$I_{j'}^{c,t}$ increasing

$\Rightarrow x^j + 1 \leq y^j$, and the order is preserved

Synthesis

Index modeling opportunities :

- Taking static informations into account :
 - Nodes characteristics : CPU speed, capacity ...
 - System characteristics : network topology
- Complex target selection strategies
 - Optimal choice : PSQ ...
 - Random probing

Impact of the control triggering

<i>Triggering policy</i>	<i>Independent control</i>	<i>Application dependent</i>
<i>Push</i>	Monotone	Monotone
<i>Pull</i>	Monotone	Non-monotone

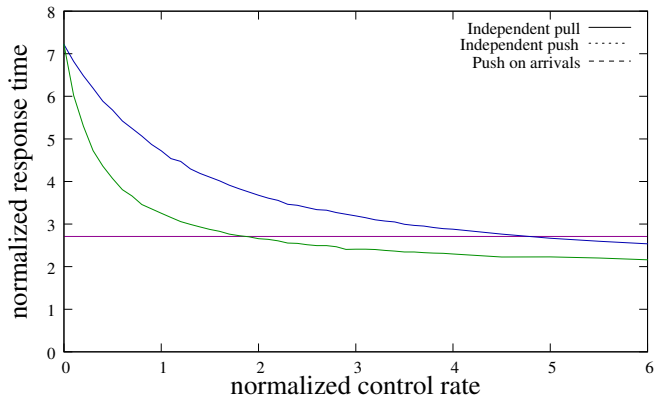
⇒ Almost monotone “Pull on completion” can be simulated with envelopes [BGV08]

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Estimation of the control rate

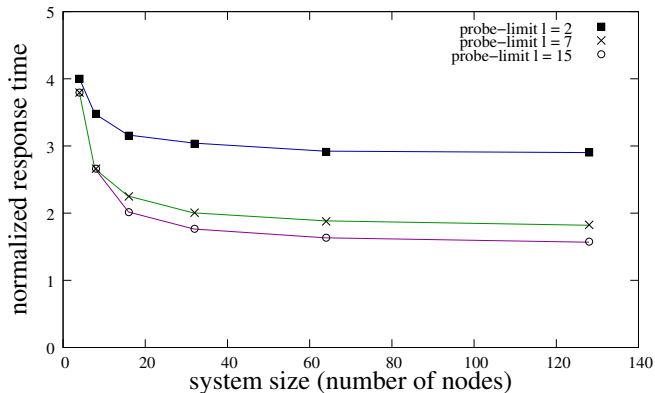
Policy : Controlled Push, Pull and Push on arrivals with random probing of 6 nodes



For a system equipped with a controller, a good operating point is to fix the control rate twice the processor speed.

Estimation of the probe-limit

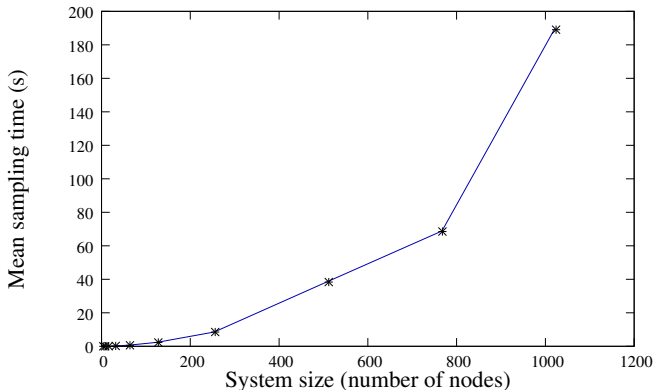
Policy : Controlled Push with random probing of 6 nodes



increasing the Probe-limit further than 7 does not provide a significant performance improvement

Scaling

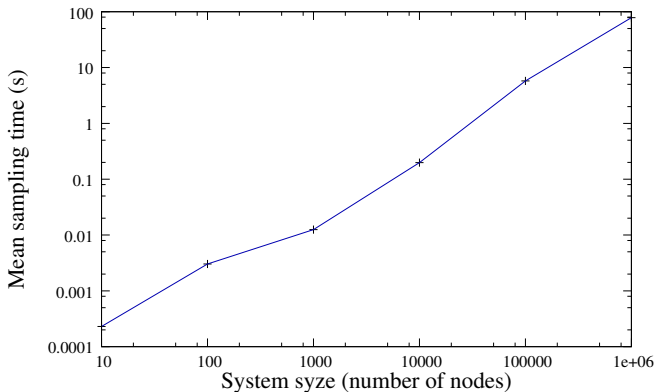
Policy : Controlled Push with random probing of 6 nodes



It is feasible to simulate complex load sharing strategies within a system of 1024 nodes.

Scaling Toward million of nodes

Policy : Threshold Push on Arrival with priority list of 8 nodes



The time to simulate such system is linear with the number of nodes



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Conclusion

A modelling framework of load sharing policies :

- complex state dependent strategies

Applications :

- Tuning of parameters
- Comparison of hierarchic work stealing strategies
- Very large scale systems

Future works :

- Comparison with mean field results

Download : <http://gforge.inria.fr/projects/psi>

INRIAForge: Perfect Simulator: Information sur un projet - Mozilla Firefox

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Résumé

PSI is a software simulator of Markov chains on large discrete state space. It samples steady state distribution in finite time by the method "coupling from the past".

- Intended Audience : End Users/Desktop
- Intended Audience : Other Audience
- Kind : Software
- License : OSI Approved ; GNU General Public License (GPL)
- Natural Language : English
- Natural Language : French
- Operating System : MacOS
- Operating System : POSIX ; Linux
- Programming Language : C
- Research center : Montbonnot
- Topic : Scientific/Engineering

Enregistré le : 24/11/2005 17:37
Taux d'activité : 30.73%
Voir les statistiques ou le rapport d'activité pour le projet.
Voir la liste des flux RSS disponibles pour ce projet.

Equipe-Projet

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[Voir les membres]
[Demander à rejoindre le projet]





Derniers fichiers publiés

Paquet	Version	Date	Remarques / Surveiller	Télécharger (download)
psi	4.4.6	March 24, 2010		Télécharger (download)
Terminé				

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